Automating Interactive Theorem Proving with Coq and Ltac

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Motivation

- Math is usually written by hand, checked by other mathematicians
- Verifying process takes time out of learning
- Computer potentially gives instant feedback

Overview

- Coq a language that verifies proofs
- Goal
 - Use Coq to help students learn proving, allowing them to check their proofs by themselves
- Problem
 - Coq is difficult to learn, and "raw Coq code" looks nothing like the original proof

"Raw Coq code" Example

specialize(H1 _ H9). intuition. intuition. rewrite H4 in H3. destruct H2 as [t]. destruct H2 as [j]. assert(Z.divide x a). unfold Z.divide. exists q. intuition. rewrite H5. ring.

A Brief Intro to Coq

- Fun Fact: Coq means rooster in French!
 Oeveloped in INRIA, France
- A basic Coq proof consists of:
 - Theorem statements
 - Tactics

• Environment

- Coq code
- List of hypotheses
- List of subgoals

Approach

• Wrote number theory proofs in Coq

Two column proof — "Raw" Coq proof Look for repetitive, messy, trivial parts Develop approach to Automate Cog Proof eliminate these parts

Approach

Proofs written:

- Well ordering principle for natural numbers
- No natural number between 0 and 1
- Infinitude of primes
 - Fundamental theorem of arithmetic
- Division algorithm
- Extended Euclidean algorithm
- Euclidean algorithm

Two Column "Math" Proof

	Theorem: there does not exist a natural number n such that $0 < n < 1$	
1	Let S be the set of all natural numbers between 0 and 1	Definition
2	Assume that S is non-empty	For sake of contradiction
3	Let r be the least element of S	Well ordering principle and 2
4	0 < r < 1	r is in S
5	0 * r < r * r < 1 * r	4 and math
6	0 < r ² < r	5 and math
7	r ² < 1	4 and 6
8	r ² is in S	6, 7 and 1
9	r ² < r	6
10	Contradiction	3 and 9
11	Since S has no least element, it is empty, and our theorem is true	Definition
12	QED	

Ugly Duckling

Theorem no_nat_between_1_and_0 : ~ exists n : nat, 0 < n < 1.

```
Proof.
 intro H.
 destruct (nat_well_ordered _ H).
  clear H.
  destruct H0.
 assert (0 * x < x * x < 1 * x).
  split; apply mult_lt_compat_r; apply H.
  assert (0 < x * x < x).
  replace 0 with (0 * x) by auto with arith.
  replace (1 * x) with x in H1 by auto with arith.
 assumption.
  assert (x * x < 1).
 apply lt_trans with x.
 apply H2.
  apply H.
 assert (0 < x * x < 1).
 split.
 apply H2.
 assumption.
  specialize (H0 _ H4).
 apply (le_not_lt _ _ H0).
 apply H2.
Qed.
```

Theorem no_nat_between_1_and_0 : ~ exists n : nat, 0 < n < 1.

~ (exists n : nat, 0 < n < 1)

Theorem no_nat_between_1_and_0 :
~ exists n : nat, 0 < n < 1.</pre>

Proof.

intro H.

H : exists n : nat, 0 < n < 1

Theorem no_nat_between_1_and_0 :
~ exists n : nat, 0 < n < 1.</pre>

Proof.

intro H.

H : exists n : nat, 0 < n < 1
x : nat
H0 : 0 < x < 1 /\ (forall a :
nat, 0 < a < 1 -> x <= a)</pre>

Theorem no_nat_between_1_and_0 :
~ exists n : nat, 0 < n < 1.</pre>

Proof.

intro H.

destruct (nat_well_ordered

H).

clear H.

x : nat H0 : 0 < x < 1 /\ (forall a : nat, 0 < a < 1 -> x <= a)</pre>

Theorem no_nat_between_1_and_0 : ~ exists n : nat, 0 < n < 1.

Proof.

intro H.

destruct (nat_well_ordered

H).

clear H.

destruct H0.



```
Theorem no_nat_between_1_and_0 :
~ exists n : nat, 0 < n < 1.
Proof.
    intro H.
    destruct (nat_well_ordered _
H).
    clear H.
    destruct H0.
    assert (0 * x < x * x < 1 *
x).</pre>
```

x : nat H : 0 < x < 1 H0 : forall a : nat, 0 < a < 1 -> x <= a 0 * x < x * x < 1 * x</pre>

subgoal 2 (ID 1226) is: False

```
Theorem no nat between 1 and 0 :
\sim exists n : nat, 0 < n < 1.
Proof.
  intro H.
  destruct (nat well_ordered
H).
  clear H.
 destruct H0.
  assert (0 * x < x * x < 1 *
X).
  split; apply mult lt compat r;
apply H.
```

```
Theorem no nat between 1 and 0 :
\sim exists n : nat, 0 < n < 1.
Proof.
  intro H.
  destruct (nat well_ordered
H).
  clear H.
  destruct H0.
  assert (0 * x < x * x < 1 *
X).
  split; apply mult lt compat r;
apply H.
  assert (0 < x * x < x).
```

```
subgoal 2 (ID 1249) is:
False
```

```
Theorem no nat between 1 and 0 :
\sim exists n : nat, 0 < n < 1.
Proof.
  intro H.
  destruct (nat well ordered
H).
  clear H.
 destruct H0.
  assert (0 * x < x * x < 1 *
X).
  split; apply mult lt compat r;
apply H.
  assert (0 < x * x < x).
  replace 0 with (0 * x) by auto
with arith.
  replace (1 * x) with x in H1
by auto with arith.
```

```
subgoal 2 (ID 1249) is:
False
```

```
Theorem no nat between 1 and 0 :
\sim exists n : nat, 0 < n < 1.
Proof.
  intro H.
  destruct (nat well_ordered
H).
  clear H.
 destruct H0.
  assert (0 * x < x * x < 1 *
X).
  split; apply mult lt compat r;
apply H.
  assert (0 < x * x < x).
  replace 0 with (0 * x) by auto
with arith.
  replace (1 * x) with x in H1
by auto with arith.
  assumption.
```

```
x : nat
H : 0 < x < 1
H0 : forall a : nat, 0 < a < 1
-> x <= a
H1 : 0 * x < x * x < 1 * x
H2 : 0 < x * x < x</pre>
```

assumption. assert (x * x < 1).

assumption.
assert (x * x < 1).
apply lt_trans with x.</pre>

x : nat H : 0 < x < 1H0 : forall a : nat, 0 < a < 1-> x <= a H1 : 0 * x < x * x < 1 * xH2 : 0 < x * x < x_____ x * x < x subgoal 2 (ID 1266) is: x < 1subgoal 3 (ID 1264) is: False

```
assumption.
assert (x * x < 1).
apply lt_trans with x.
apply H2.
apply H.</pre>
```

x : nat H : 0 < x < 1 H0 : forall a : nat, 0 < a < 1 -> x <= a H1 : 0 * x < x * x < 1 * x H2 : 0 < x * x < x H3 : x * x < 1 False

assumption.
assert (x * x < 1).
apply lt_trans with x.
apply H2.
apply H.
assert (0 < x * x < 1).</pre>

x : nat H : 0 < x < 1H0 : forall a : nat, 0 < a < 1-> x <= a H1 : 0 * x < x * x < 1 * xH2 : 0 < x * x < xH3 : x * x < 10 < x * x < 1subgoal 2 (ID 1272) is: False

```
assumption.
assert (x * x < 1).
apply lt_trans with x.
apply H2.
apply H.
assert (0 < x * x < 1).
split.</pre>
```

```
x : nat
 H : 0 < x < 1
  H0 : forall a : nat, 0 < a < 1
-> x <= a
  H1 : 0 * x < x * x < 1 * x
  H2 : 0 < x * x < x
  H3 : x * x < 1
  0 < x * x
subgoal 2 (ID 1275) is:
x * x < 1
subgoal 3 (ID 1272) is:
False
```

```
assumption.
assert (x * x < 1).
apply lt_trans with x.
apply H2.
apply H.
assert (0 < x * x < 1).
split.
apply H2.
assumption.
```

```
x : nat
H : 0 < x < 1
H0 : forall a : nat, 0 < a < 1
-> x <= a
H1 : 0 * x < x * x < 1 * x
H2 : 0 < x * x < x
H3 : x * x < 1
H4 : 0 < x * x < 1
False
```

assumption. assert (x * x < 1). apply lt_trans with x. apply H2. apply H. assert (0 < x * x < 1). split. apply H2. assumption. specialize (H0 H4). x : nat H : 0 < x < 1 H0 : x <= x * x H1 : 0 * x < x * x < 1 * x H2 : 0 < x * x < x H3 : x * x < 1 H4 : 0 < x * x < 1</pre>

assumption. assert (x * x < 1). apply lt_trans with x. apply H2. apply H. assert (0 < x * x < 1). split. apply H2. assumption. specialize (H0 _ H4). apply (le_not_lt _ H0). x : nat H : 0 < x < 1 H0 : x <= x * x H1 : 0 * x < x * x < 1 * x H2 : 0 < x * x < x H3 : x * x < 1 H4 : 0 < x * x < 1 x * x < x</pre>

assumption. assert (x * x < 1). apply lt_trans with x. apply H2. apply H. assert (0 < x * x < 1). split. apply H2. assumption. specialize (H0 _ H4). apply (le_not_lt _ H0). apply H2. No more subgoals. (dependent evars:)

```
assumption.
assert (x * x < 1).
apply lt_trans with x.
apply H2.
apply H.
assert (0 < x * x < 1).
split.
apply H2.
assumption.
specialize (H0 _ H4).
apply (le_not_lt _ H0).
apply H2.
```

no_nat_between_1_and_0 is
defined

● ○ ○ □ no_nat_betw	veen_0_and_1.v
👀 👀 📟 🗶 🔺 🕨 🗶 🛏 💓 🖀 🛹 🐧 🐖 🖧 🗧	9
<pre>Theorem no_nat_between_1_and_0 : ~ exists n : nat, 0 < n < 1. Proof. intro H. destruct (nat_well_ordered _ H). clear H. destruct H0. assert (0 * x < x * x < 1 * x). split; apply mult_lt_compat_r; apply H. assert (0 < x * x < x). replace 0 with (0 * x) by auto with arith. replace (1 * x) with x in H1 by auto with arith. assumption. assert (x * x < 1). apply lt_trans with x. apply H2. apply H. assert (0 < x * x < 1). split. apply H2. assumption. specialize (H0 _ H4). cmply (lo mot lt = H0)</pre>	-:%%- *goals* All L1 (Coq Goals)
apply (le_not_lt H0). apply H2.	no_nat_between_1_and_0 is defined
Qed.	
U:**- no_nat_between_0_and_1.v 57% L106 (Coq Script(0-) Hole	s -:%%- *response* All L1 (Coq Response)

Proof Automation

- Noticed repeated, tedious code
- Goal less cluttered code, easier to write
- Write tactics that automate trivial steps
- Three tactics written
 - Zwop
 - math
 - \circ simplify

"Zwop" Tactic

- Applies the well-ordering principle (WOP) to a set in Coq
- WOP is a key part of many number theory proofs

```
Ltac Zwop A Adef :=
   pose(A := Adef);
   destruct (bounded_Z_well_ordered A).
```

"math" Tactic

• Tries to simplify ring inequalities and uses "intuition" tactic

```
Ltac math :=

intuition;

repeat match goal with

I [ H : ?T, H' : ?T' I- _ ] =>

match type of T with

I Prop => match type of T' with

I Prop => assert (T /\ T') by tauto; clear H H'

end

end

end;

repeat match goal with

I [ H : context[?a < ?b] I- _ ] => progress ring_simplify a b in H;

progress apply lt_trans with a in H; progress apply lt_trans with b in H

end;

intuition.
```

"simplify" Tactic

• Tries many general, useful tactics where applicable

• Only uses them if they advance the proof

```
Ltac simplify_step v :=
  repeat (try match goal with
           | [ |- _ ∧ _] ⇒ split
           | [H: \_ \land \_ |- \_] \Rightarrow destruct H
           [ [ - _ = _] => ring
           I [ H : exists _, _ I- _] => let v' := fresh v in destruct H as [v']
           | [|- exists _, _] \Rightarrow eauto with arith
           end:
  intuition;
  simpl in *;
  unfold Z.divide in *).
Ltac simplify' v :=
  simplify_step v;
  try now (subst; simplify_step v).
Tactic Notation "simplify" := let v := fresh "v" in simplify' v.
Tactic Notation "simplify" ident(v) := simplify' v.
```

The Power of "simplify"

From this...

```
Lemma remainder_not_0 : forall a b, remainder a b = 0 -> a <> 0 -> b <> 0.
Proof.
    intros.
    apply NNPP; intro.
    assert (b = 0); intuition.
    subst.
    unfold remainder in H; intuition.
Qed.
```

The Power of "simplify"

To this...

Proof. simplify. Qed.

Lists From Sets

- Synthesizes list of elements from finite set
- Allows us to compute operations on the set

```
Lemma finite_set_list : forall T (E : Ensemble T), Finite T E ->
exists L : list T, forall t : T, List.In t L <-> Ensembles.In _ E t.
```

Assertion Technique

- Coq uses backwards proving
- Forward proving is more natural
- Used assert tactic to do so
- Three steps
 - Write out two column proof
 - Convert the left column into assertions in the proof
 - Prove each assertion using the ones before it

Basic Frameworks

Theorem theorem_name : [theorem_statement]. Proof.

intros.
assert([step 1]).
 [code used to prove step 1]
assert([step 2]).
 [code used to prove step 2]

Qed.

Beautiful Swan

```
Theorem no_nat_between_1_and_0 : ~ exists n : nat, 0 < n < 1.
Proof.
  intro.
  pose (U := (fun n => 0 < n < 1)).
  destruct (nat_well_ordered U) as [r]; [ auto | ].
  unfold U in HO.
  assert (0*r < r*r < 1*r) by (apply mult_double_lt_compat_r; math).
  assert (0 < r*r < r) by math.
  assert (r*r < 1) by (apply lt_trans with r; math).
  assert (r*r \in U) by (unfold U; math).
  assert (r*r < r) by math.
  assert (r <= r*r) by math.
  omega.
Oed.
```

Conclusion

- Our findings have educational implications
 - Coq proofs now resemble mathematical proofs more
 - More readable
 - Proofs are easier to write
 - Don't have to do everything out
 - Easier for students to use Coq to learn to write proofs and validate their proofs

Future

- Interpreter
 - Inputs higher level language, outputs Coq code
- Doing out more Number Theory proofs in Coq for more automation

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