# Counting the Involutions of the Symmetric Group

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### **Background**



### **Robinson-Schensted** Algorithm



Viennot 's 03 | Geometric Construction

# number involutions of $S_n$ $\Leftrightarrow$ std tableaux with n elements

### **Definitions**

### Symmetric Group

A collection of the permutations of {1,..., n} that acts as a group with composition of permutations being the group operation

There are n! elements

Ex.

We write elements as two rows  $\rightarrow$  1 2 3  $\leftarrow$  omit the first row for simplicity 1 3 2

132 is an element of S<sub>3</sub> where

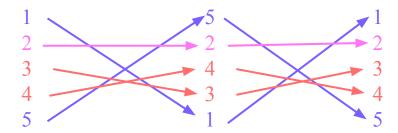
- $1 \rightarrow 1$
- $2 \rightarrow 3$
- $3 \rightarrow 2$

### **Definitions**

### Involution

For S<sub>n</sub>, an element is an involution if applying the permutation twice maps every element to itself

Ex. 52431 in S<sub>5</sub>



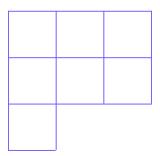


### **Partition**

# Unordered decomposition of a positive integer into positive integer parts

eg. 
$$7 = 3 + 3 + 1$$

### Ferrers Diagram



*shape*  $\lambda = (3,3,1)$ 

### **Tableau**

5	3	1
2	4	6
7		

### **Standard Tableau**

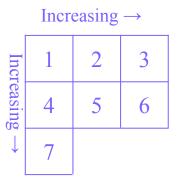
Increasing  $\rightarrow$ 

Incr	1	2	3
ncreasing	4	5	6
<b>J</b>	7		

1	3	5
2	4	6
7		

1	3	4
5	2	7
6		•

### **Standard Tableau**



1	3	5
2	4	6
7		

1	3	4
5	2	7
6		

standard

standard

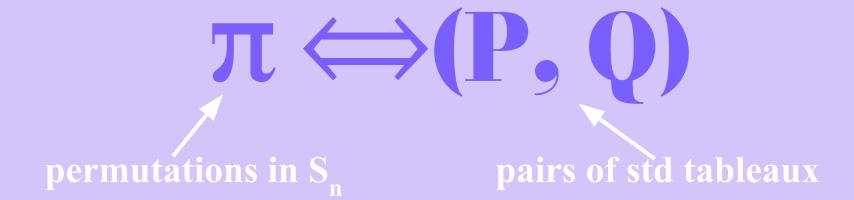
not standard

# number involutions of $S_n$ $\Leftrightarrow$ std tableaux with n elements



# R-S Algorithm





### Forward Correspondence

### P "insertion tableau"

```
given \pi = x_1 \overline{x_2 \dots x_n}
for i in 1,..., n {
      R = 1
      X = X_i
      while x < some element in row R {
            y = min element of row R st y > x
            replace y by x
            set x:= y
            R++
      append x to the end of row R
```

### Q "recording tableau"

```
given \pi = x_1 x_2 ... x_n
for i in 1,..., n {
 set the cell in which insertion terminates to n }
```

P

4

# Q

#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_1
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x to the end of row R
```

```
for i in 1,..., 5 {
    set the cell in which insertion terminates to n
}
```

inserted: 4

DISPLACED: n/a

P

4

Q

1

### P TABLEAU

```
for i in 1,..., 5 {
     R = 1
     X = X_1
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x to the end of row R
```

### **Q TABLEAU**

```
for i in 1,..., 5 {
```

set the cell in which insertion terminates to  $m{1}$ 

### INSERTING: 1 DISPLACED: 4

P

1

Q

1

```
P TABLEAU
for i in 1,..., 5 {
     R = 1
     X = X_2
     while x < some element in row 1 {
           y = min element of row 1 st y > x = 4
           replace y by x
           set x:= y
           R++
     append x to the end of row R
Q TABLEAU
for i in 1,..., 5 {
     set the cell in which insertion terminates to n
```

### INSERTING: 4 DISPLACED: n/a

P

1 4

Q

1

### P TABLEAU

```
for i in 1,..., 5 {
     R = 1
     X = X_2
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x = 4 to the end of row 2
```

```
for i in 1,..., 5 {
    set the cell in which insertion terminates to n
}
```

### inserted: 1 DISPLACED: 4

P

1 4

Q

2

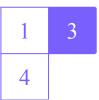
#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_2
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x to the end of row R
```

```
for i in 1,..., 5 {
    set the cell in which insertion terminates to 2
}
```

### INSERTING: 3 DISPLACED: n/a

P



Q

```
2
```

#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_3
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append 3 to the end of row 1
```

```
for i in 1,..., 5 {
    set the cell in which insertion terminates to n
}
```

### inserted: 3 DISPLACED: n/a

P

1	3
4	

Q

```
1 3
```

#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_3
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x to the end of row R
```

```
for i in 1,..., 5 {
    set the cell in which insertion terminates to n
}
```

### INSERTING: 2 DISPLACED: 3

P

1	2
4	

Q

```
1 3
```

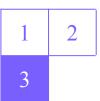
#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_{\Lambda}
      while x < some element in row R {
            y = min element of row R st y > x
            replace y by x
            set x:= y
            R++
      append x to the end of row R
```

```
for i in 1,..., 5 {
     set the cell in which insertion terminates to n
}
```

### INSERTING: 3 DISPLACED: 4

P



Q

```
1 3
```

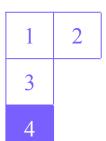
#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_{\Lambda}
      while x < some element in row R {
            y = min element of row R st y > x
            replace y by x
            set x:= y
            R++
      append x to the end of row R
```

```
for i in 1,..., 5 {
     set the cell in which insertion terminates to n
}
```

### INSERTING: 4 DISPLACED: n/a

P



Q

```
1 3
```

#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_{\Lambda}
     while x < some element in row R {
            y = min element of row R st y > x
            replace y by x
            set x:= y
            R++
     append x to the end of row R
```

```
for i in 1,..., 5 {
     set the cell in which insertion terminates to n
}
```

inserted: 2

DISPLACED: 3, 4

P

1	2
3	
4	

Q

```
1 3
2
```

#### **P TABLEAU**

```
for i in 1,..., 5 {
      R = 1
     X = X_{\Lambda}
      while x < some element in row R {
            y = min element of row R st y > x
            replace y by x
            set x:= y
            R++
      append x to the end of row R
```

```
for i in 1,..., 5 {
    set the cell in which insertion terminates to n
}
```

### INSERTING: 5 DISPLACED: n/a

P

1	2	5
3		
4		

Q

```
1 3
2
4
```

#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_5
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x to the end of row R
```

```
for i in 1,..., 5 {
    set the cell in which insertion terminates to n
}
```

### inserted: 5 DISPLACED: n/a

P

1	2	5
3		
4		

Q

```
1 3 5
2 4
```

#### **P TABLEAU**

```
for i in 1,..., 5 {
     R = 1
     X = X_5
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x to the end of row R
Done!
Q TABLEAU
for i in 1,..., 5 {
     set the cell in which insertion terminates to n
```

### INSERTING: // DISPLACED: //

P

1	2	5
3		
4		

Q

```
1 3 5
2
4
```

#### **P TABLEAU**

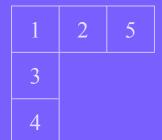
```
for i in 1,..., 5 {
     R = 1
     X = X_i
     while x < some element in row R {
           y = min element of row R st y > x
           replace y by x
           set x:= y
           R++
     append x to the end of row R
Done!
Q TABLEAU
for i in 1,..., 5 {
     set the cell in which insertion terminates to n
```

# $\pi \rightarrow (P, 0)$

$$(41325) \longrightarrow \begin{pmatrix} \begin{vmatrix} 1 & 2 & 5 \\ 3 & & \\ 4 & & 4 \end{pmatrix}$$

# $\pi \leftarrow (P, Q)$

# P





1	3	5
2		
4		

### From Q:

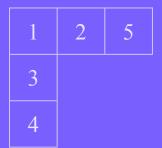
take *largest cell* and find its coordinates (*i*, *j*) and content *index* 

#### In P:

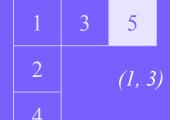
```
take element in (i, j) and call it x
let row be row R
while row R isn't first row:
    find the largest element y of row R-1 less than x
    replace y by x
    set x := y
    R--
```



# P







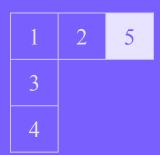
### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

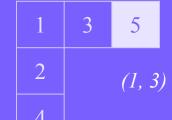
#### In P:

```
take element in (i, j) and call it x
let row be row R
while row R isn't first row:
    find the largest element y of row R-1 less than x
    replace y by x
    set x := y
    R--
```

# P







### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

#### In P:

take element in *(i, j)* and call it *x* let row be row R

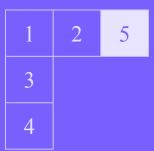
while row R isn't first row:

find the largest element y of row R-1 less than x replace y by x

set *x := y* 

R---

# P





1	3	5
2		(1, 3)
4		

### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

#### In P:

```
take element in (i, j) and call it x

let row be row R

while row R isn't first row:

find the largest element y of row R-1 less than x

replace y by x

set x := y

R--
```

## P



3

4



```
1 3
```

2

4

### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

#### In P:

```
take element in (i, j) and call it x
let row be row R
while row R isn't first row:
    find the largest element y of row R-1 less than x
    replace y by x
    set x := y
    R--
```

 $\begin{array}{c|cccc}
 & 1 & 2 \\
\hline
 & 3 & \\
\hline
 & 4 & \\
\end{array}$ 



### From Q:

take *largest cell* and find its coordinates (i, j) and content idx

#### In P:

```
take element in (i, j) and call it x
let row be row R
while row R isn't first row:
    find the largest element y of row R-1 less than x
    replace y by x
    set x := y
    R--
```

### 3 4 1 3 2

(3, 1)

### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

#### In P:

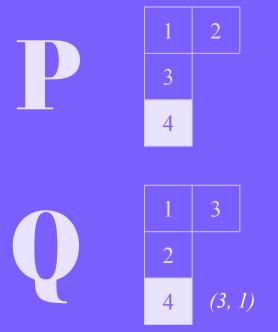
take element in *(i, j)* and call it *x* let row be row R

while row R isn't first row:

find the largest element y of row R-1 less than x replace y by x

set *x := y* 

R---



### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

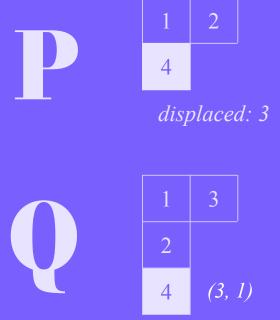
#### In P:

take element in (i, j) and call it x let row be row R

while row R isn't first row:

find the largest element y of row R-1 less than x replace y by x set x := y

# $\pi = 5$



#### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

#### In P:

```
take element in (i, j) and call it x

let row be row R

while row R isn't first row:

find the largest element y of row R-1 less than x

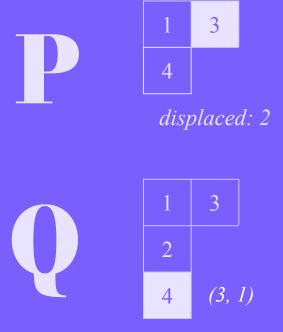
replace y by x

set x := y

R--
```

when row one reached, x is the index idx element of the permutation

# $\pi =$ 5



#### From Q:

take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

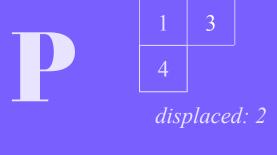
#### In P:

```
take element in (i, j) and call it x let row be row R while row R isn't first row:
```

```
find the largest element y of row R-1 less than x replace y by x set x := y R--
```

when row one reached, *x* is the index *idx* element of the permutation

## $\pi = 25$





#### From Q:

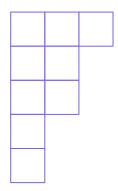
take *largest cell* and find its coordinates (*i*, *j*) and content *idx* 

#### In P:

```
take element in (i, j) and call it x
let row be row R
while row R isn't first row:
    find the largest element y of row R-1 less than x
    replace y by x
    set x := y
    R--
```

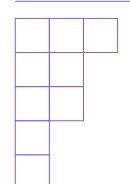
when row one reached, x is the index idx element of the permutation

# π <del>= (P, Q)</del>

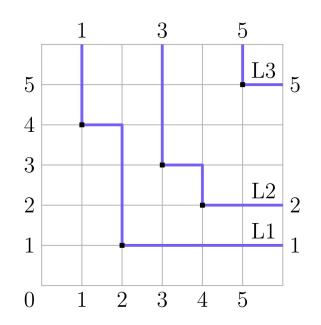


# Viennot's Construction

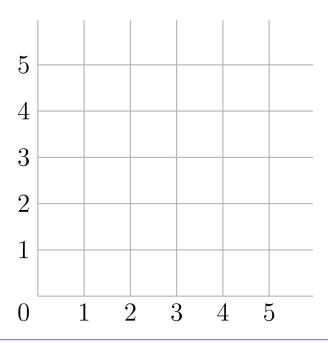
Correspondence to a Shadow Line Diagram



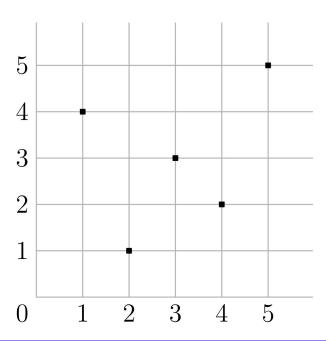


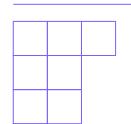




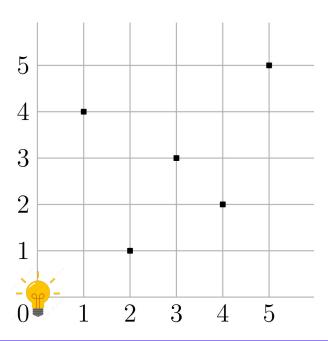


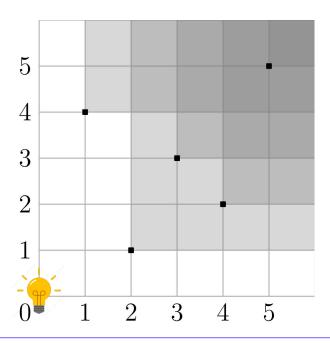
1) given  $\pi = x_1 x_2 \dots x_n$  represent  $x_i$  by a box with coordinates  $(i, x_i)$ 

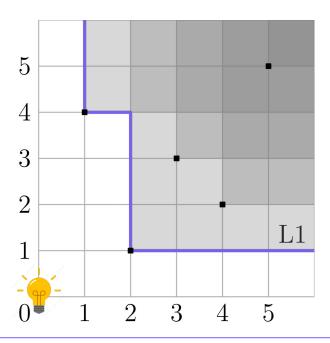


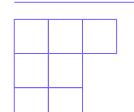


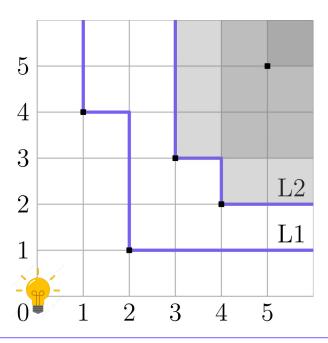
2) put a light at the origin

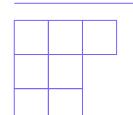


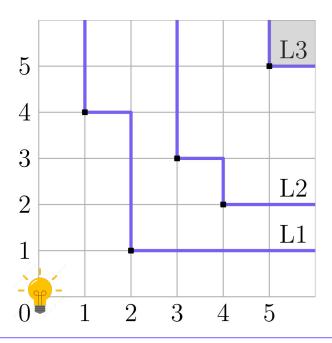


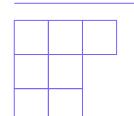




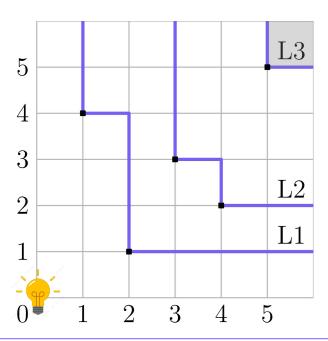




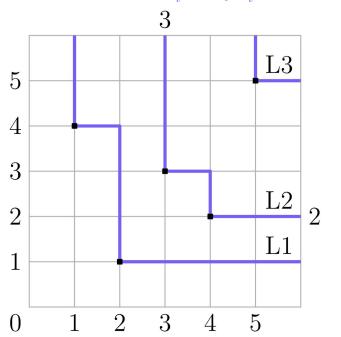




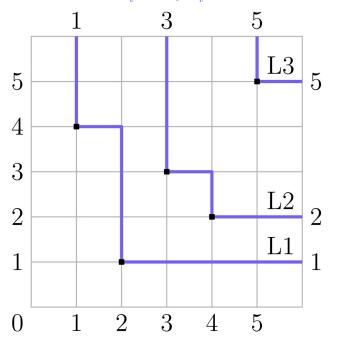
4) label shadow line coordinates –  $\,x_{L_i}$  and  $\,y_{L_i}$ 

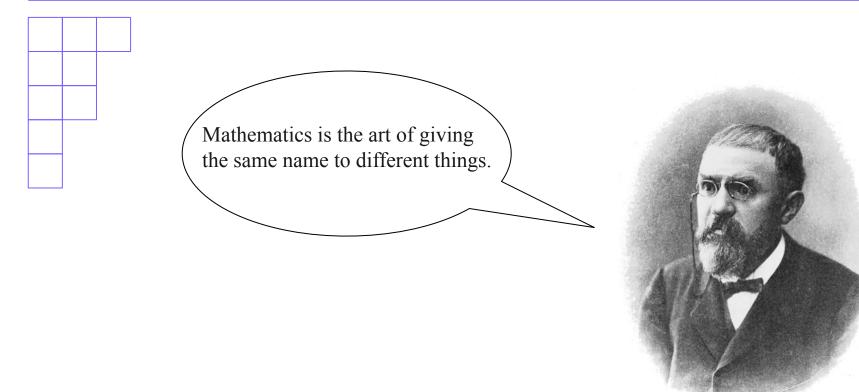


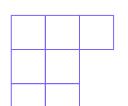
4) label shadow line coordinates –  $x_{L_i}$  and  $y_{L_i}$ 



4) mark shadow line labels –  $\,x_{L_i}$  and  $\,y_{L_i}$ 



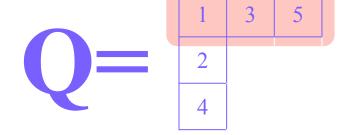




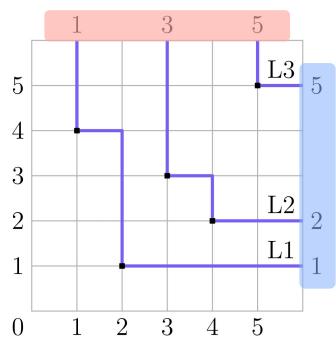
#### The secret correspondence $\Leftrightarrow$ (for $\pi = 41325$ )

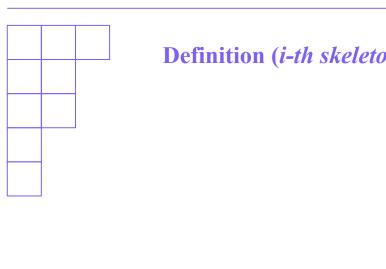


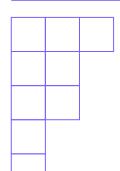
R-S algorithm:



Viennot's construction:

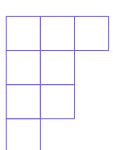




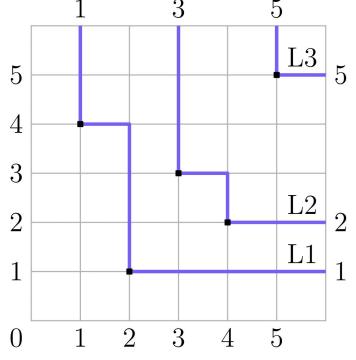


$$\pi^{(1)} = \pi \text{ and }$$

where  $(k_1, l_1), ..., (k_m, l_m)$  are the northeast corners of the shadow diagram of  $\pi^{(i-1)}$  listed in lexicographic order. The shadow lines of  $\pi^{(i)}$  are denoted by  $L_i^{(i)}$ .



#### **Definition** (*i-th skeleton of* $\pi$ ): $\pi^{(1)} = \pi$ and

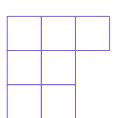


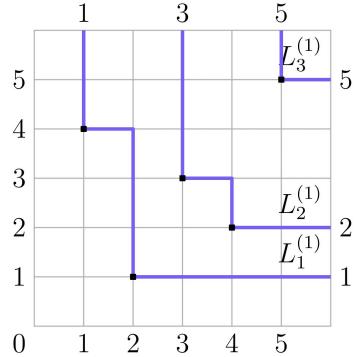
$$\pi^{(1)} = \pi \text{ and }$$

where  $(k_1, l_1), ..., (k_m, l_m)$  are the northeast corners of the shadow diagram of  $\pi^{(i-1)}$  listed in lexicographic order. The shadow lines of  $\pi^{(i)}$  are denoted by  $L_i^{(i)}$ .

# EXAMPLE

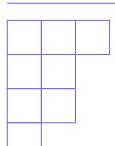
$$\pi = 41325$$

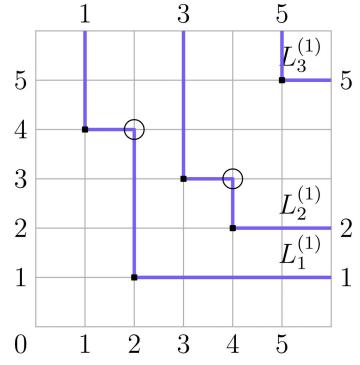




$$\pi^{(1)} = \frac{1}{4} \quad \frac{2}{1} \quad \frac{3}{3} \quad \frac{4}{3} \quad \frac{5}{5}$$

EXAMPLE 
$$\pi = 41325$$

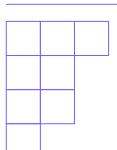


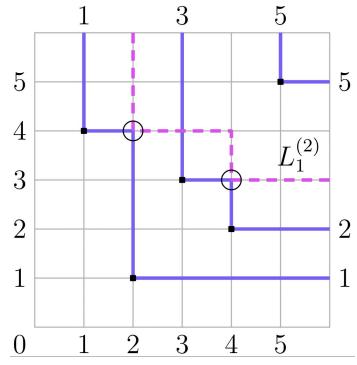


$$\pi^{(1)} = \frac{1}{4} \quad \frac{2}{1} \quad \frac{3}{3} \quad \frac{4}{2} \quad \frac{5}{5}$$

$$\left(\pi^{(2)} = \frac{2}{4} \quad \frac{4}{3}\right)$$

EXAMPLE 
$$\pi = 41325$$

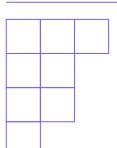


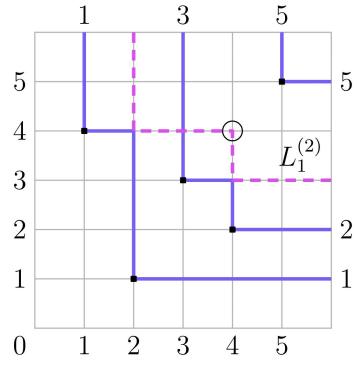


$$\pi^{(1)} = \frac{1}{4} \quad \frac{2}{1} \quad \frac{3}{3} \quad \frac{4}{2} \quad \frac{5}{5}$$

$$\left(\pi^{(2)} = \begin{array}{cc} 2 & 4 \\ 4 & 3 \end{array}\right)$$

EXAMPLE 
$$\pi = 41325$$



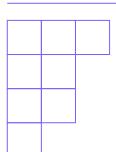


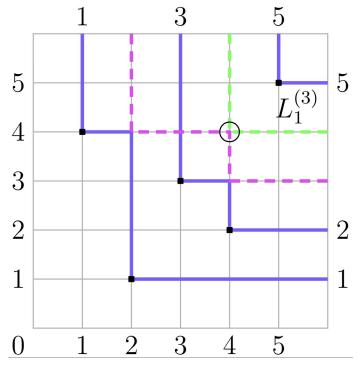
$$\pi^{(1)} = \frac{1}{4} \quad \frac{2}{1} \quad \frac{3}{3} \quad \frac{4}{2} \quad \frac{5}{5}$$

$$\left(\pi^{(2)} = \frac{2}{4} \quad \frac{4}{3}\right)$$

$$\pi^{(3)} = \frac{4}{4}$$

$$\pi = 41325$$





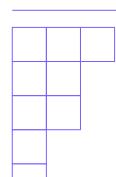
$$\pi^{(1)} = \frac{1}{4} \quad \frac{2}{1} \quad \frac{3}{3} \quad \frac{4}{2} \quad \frac{5}{5}$$

$$\pi^{(2)} = \frac{2}{4} \quad \frac{4}{3}$$

$$\pi^{(3)} = \frac{4}{4}$$

**EXAMPLE** 

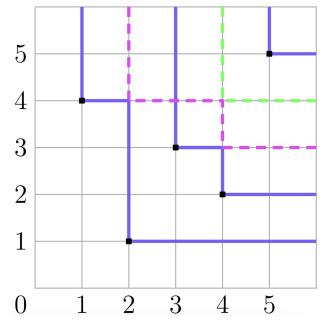
$$\pi = 41325$$



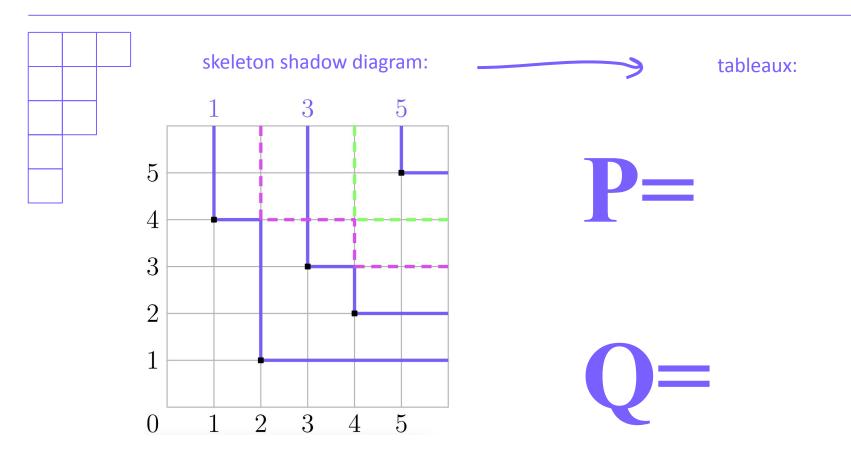
#### skeleton shadow diagram:

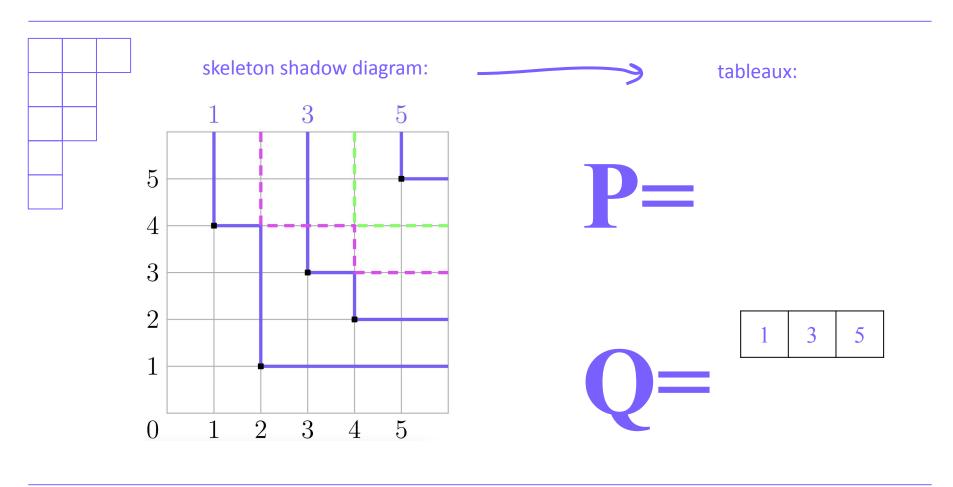


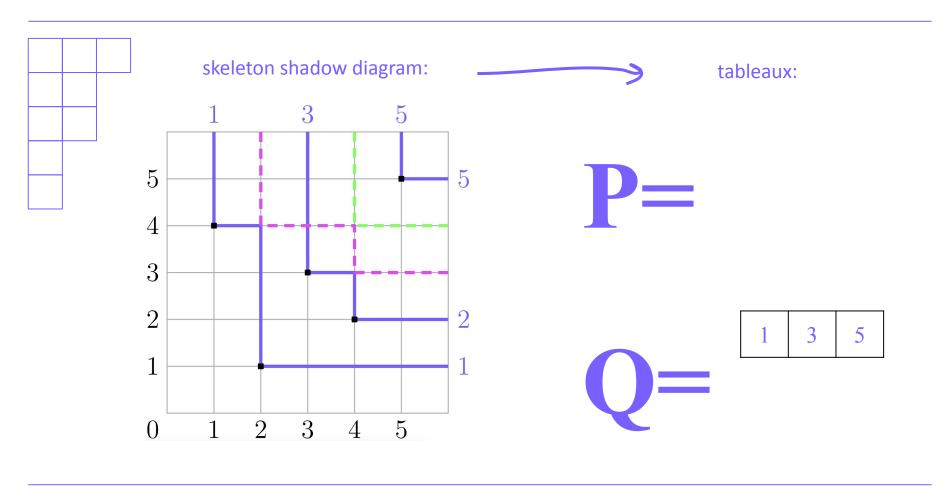
tableaux:

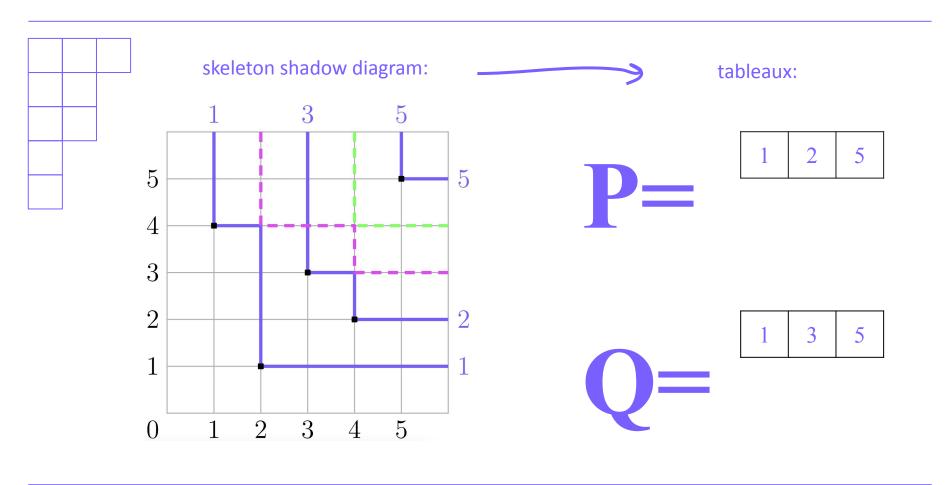


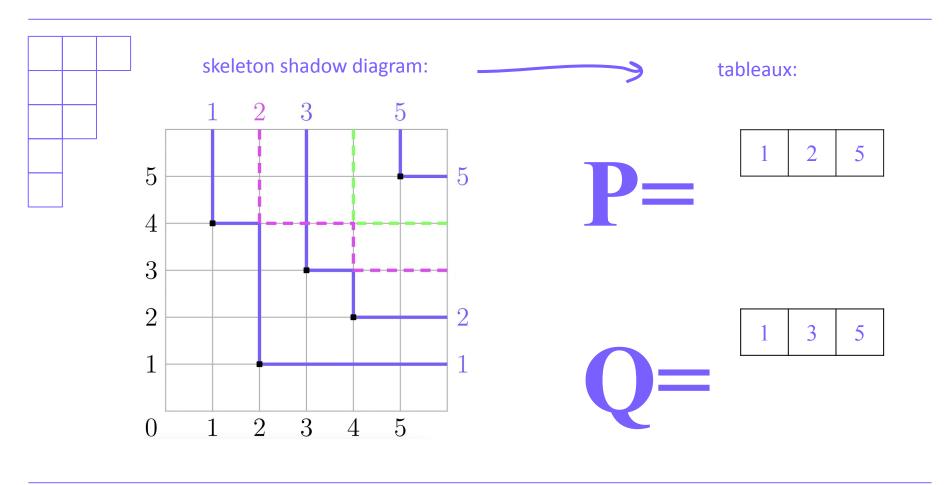


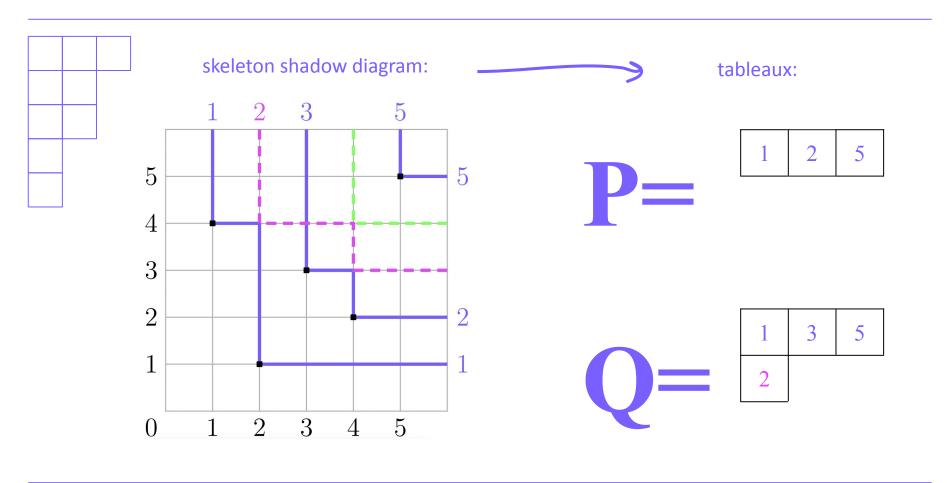


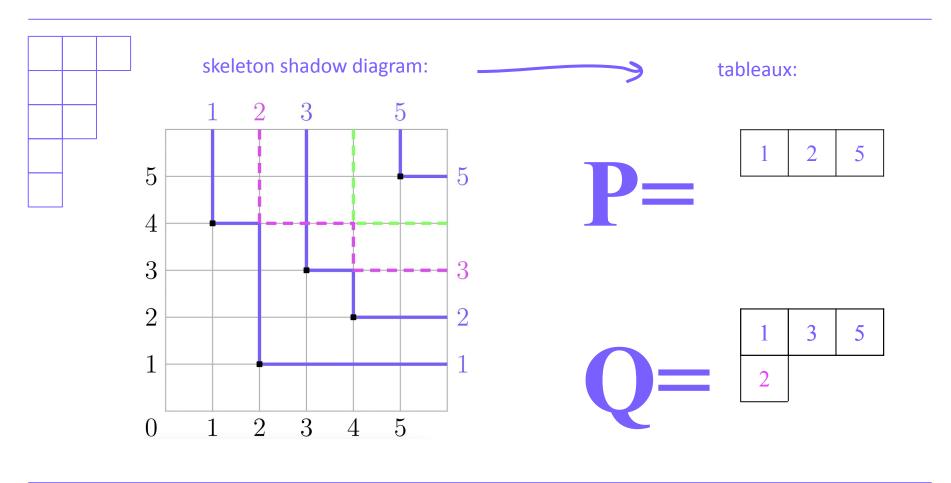


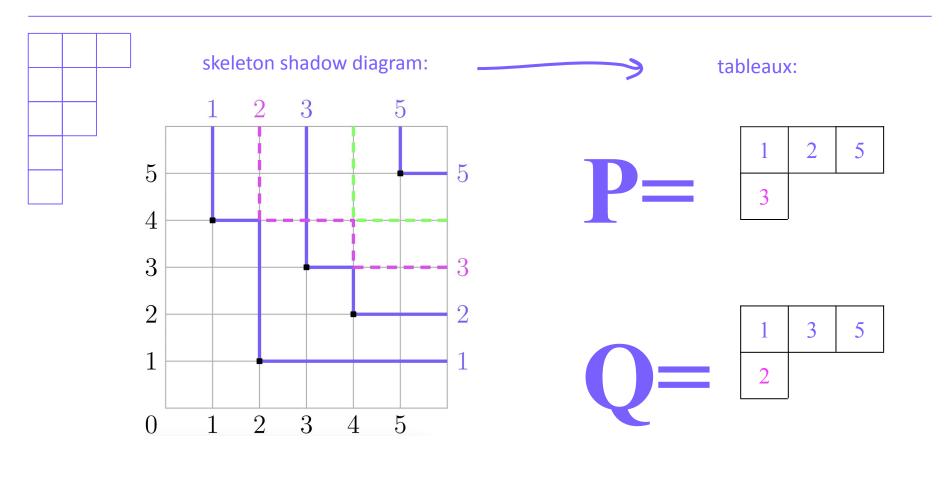


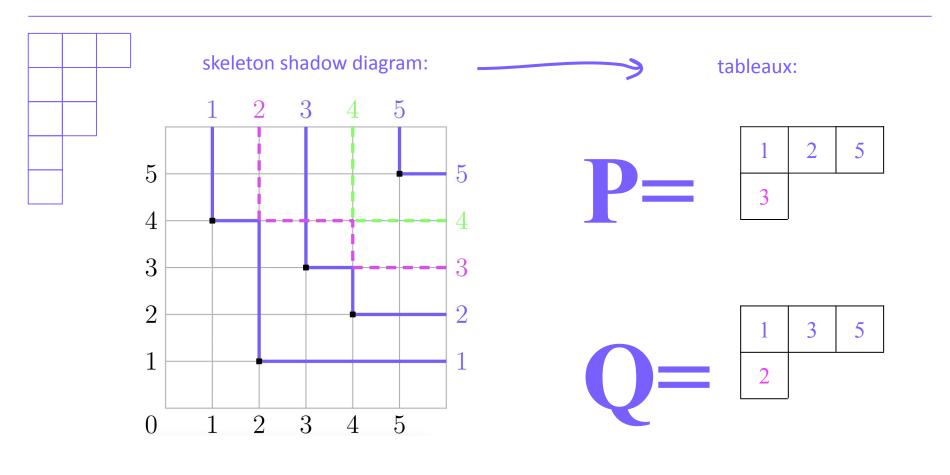


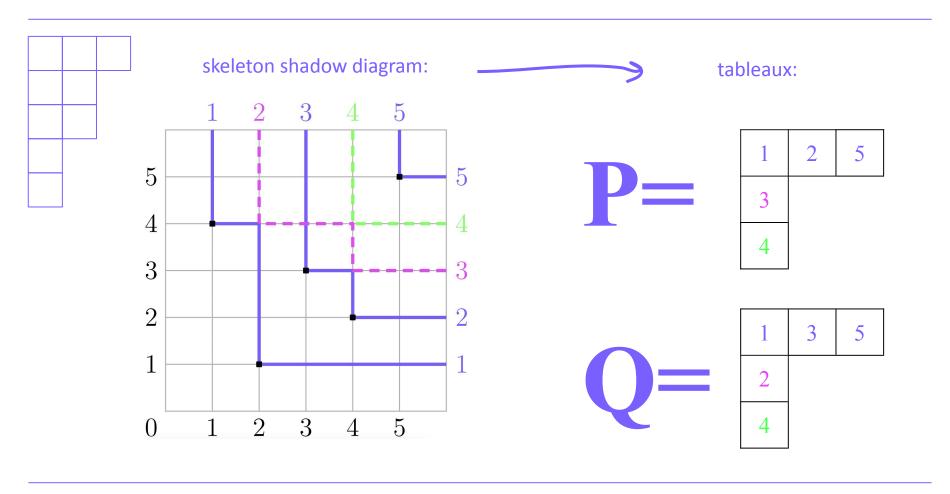


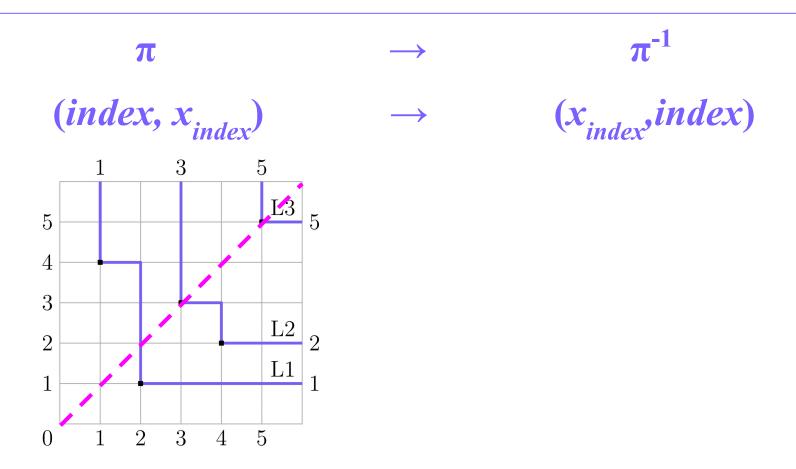


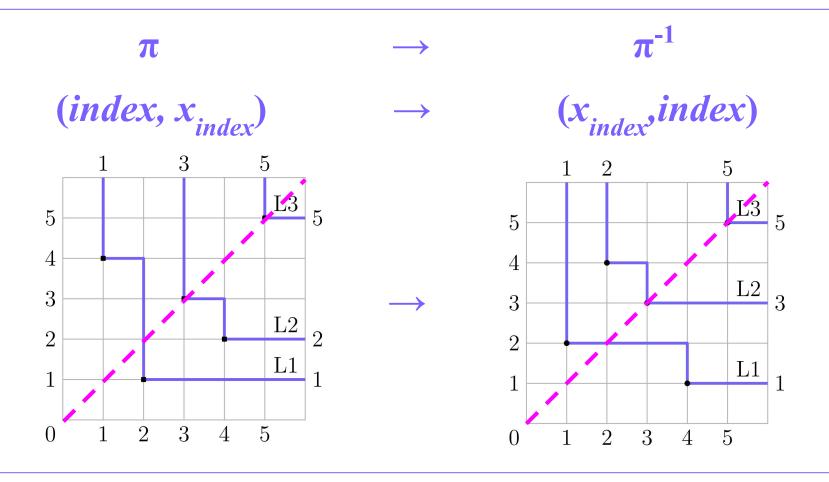


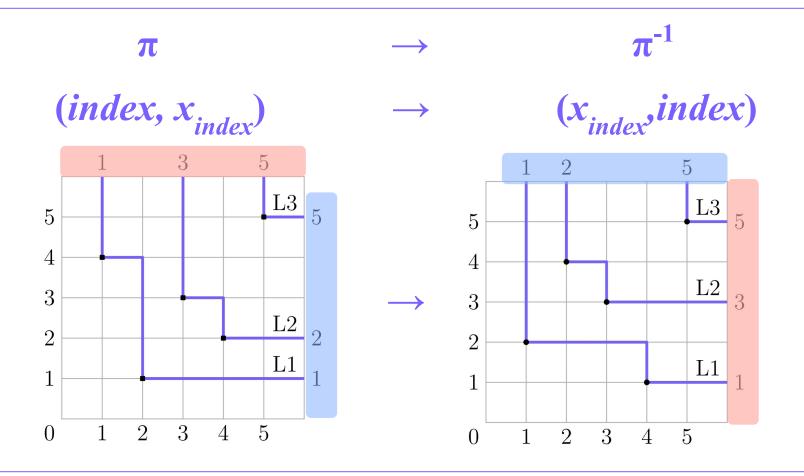


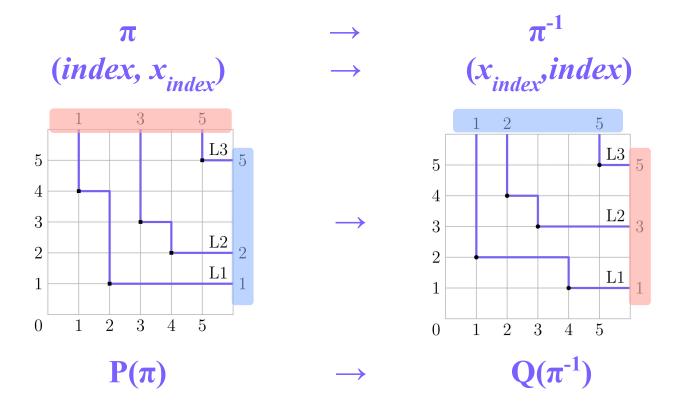










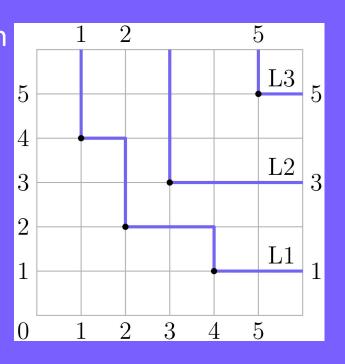


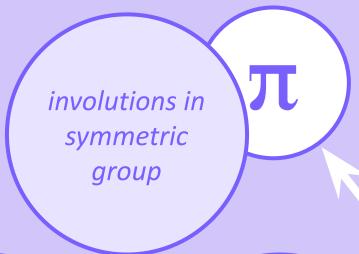
for *involution*,  $\pi = \pi^{-1}$  $P(\pi) = Q(\pi^{-1}) \Rightarrow P(\pi) = Q(\pi)$  for  $\pi$  involution

then, there is a bijection from *involutions* to pairs of identical standard tableaux ( $P(\pi)$ ,  $P(\pi)$ )

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Note: involutions correspond to shadow diagrams that are *symmetric* about y=x

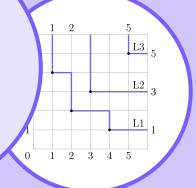




pairs of identical standard tableaux P, Q

(P, Q)

shadow diagrams symmetric about y=x





# involutions

equals

# std tableau

# Example with $S_4$

Standard Tableaux with 4 elements: 10

1	2	3	4
---	---	---	---

3	4	

1	
2	
3	
1	

1	2
3	
1	

1	3
2	
4	

1	4
2	
3	

1	2
3	4

1	3
2	4

1	2	3
4		

1	3	4
2		

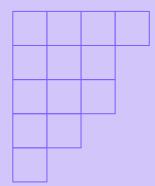
1	2	4
3		

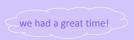
## Example with $S_4$

Involutions in  $S_4$ : 10

1234 1243 1432 1324 2134 3214 4231

2143 3412 4321







# Thank you, PRIMES and parents!

We would like to thank our mentor, Serina Hu, PRIMES coordinators Prof. Pavel Etingof, Dr. Slava Gerovitch, and Dr. Tanya Khovanova, everybody behind the PRIMES program, as well as our parents!